

# FUN WITH INFINITY

The Hilbert Hotel is very large. It has an infinite number of rooms - one for each natural number.

It is possible (even in theory) for the Hilbert Hotel to be full? What if full means:

1) Each room is occupied

or

2) If anyone else came, the hotel could not accommodate them

# SET THEORY

Monday, 29 November

# FORMAL SEMANTICS

- An interpretation (or structure in LPL chap 18) has:
  - A set of objects as the domain
  - A set of objects for each one place predicate (the set of things that satisfy the predicate)
  - A set of ordered pairs for each two place predicate [n-tuples for the n-place predicates]
  - A function which maps names into objects in the domain
  - Assigns a function (domain to domain) for each function in the language

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Takes care of P1, P2



$$\text{Domain: } \{1,2,3,4\}$$

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Added for P3





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Of course there is this mathematical theory - set theory - written in FOL where there are axioms giving the formal definition of set. In set theory there is one 2-place predicate  $E(x,y)$  and  $a \in b$  is just shorthand for  $E(a,b)$

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To formalize this we have the first axiom: Extensionality

$$\forall x \forall y [\forall z (z \in x \leftrightarrow z \in y) \rightarrow x = y]$$

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$\{x \mid x \text{ is a person in this room}\}$

$\{\{x,y\} \mid x,y \text{ live in Ithaca and } x,y \text{ are married to each other}\}$

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This leads to some paradoxes e.g. The set of all truths

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Really, this is an axiom scheme with an infinite number of axioms:

$\exists y \forall x [x \in y \leftrightarrow x \neq x]$  - the empty set

$\exists y \forall x [x \in y \leftrightarrow x = 3]$  - the singleton set  $\{3\}$

$\exists y \forall x [x \in y \leftrightarrow (x \in a \vee x \in b)]$  - the union of  $a, b$

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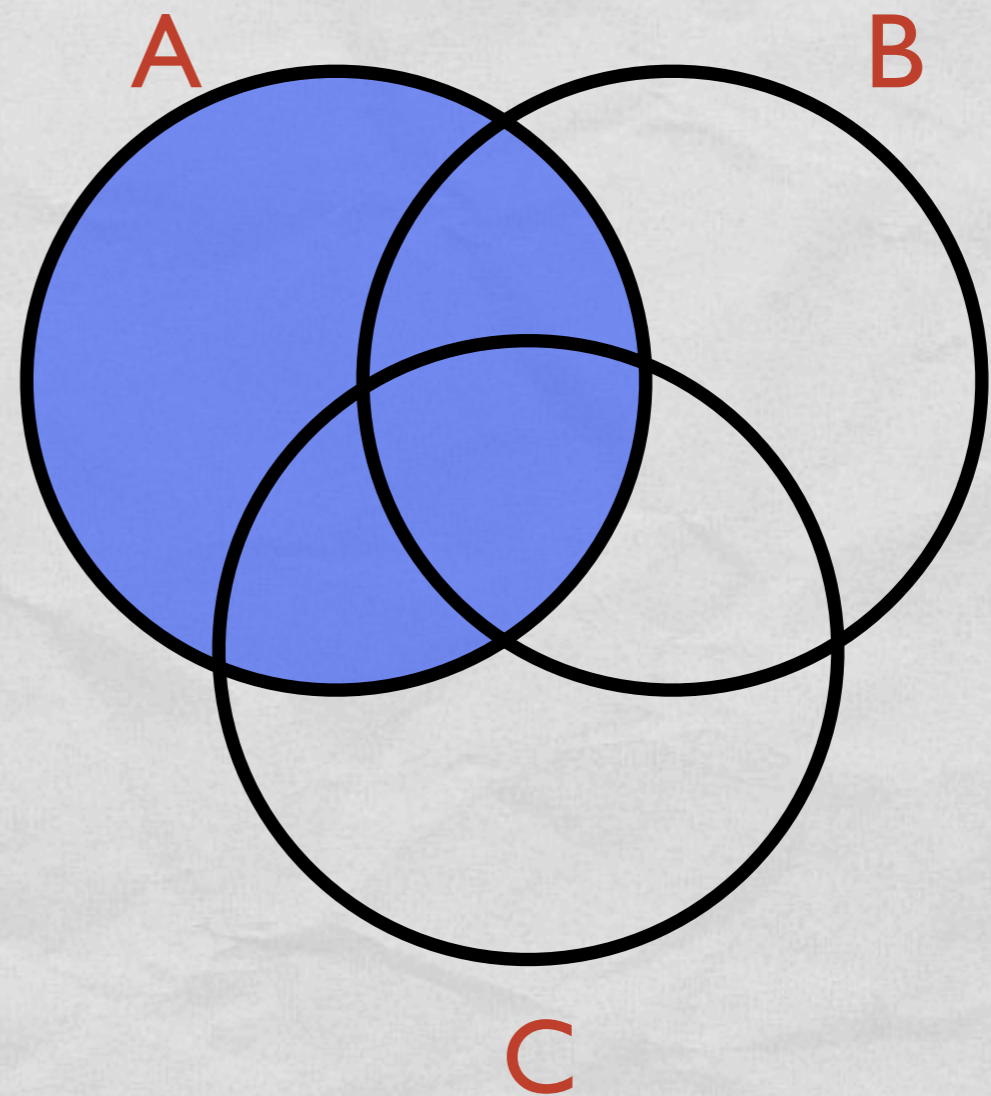
$$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$



# VENN DIAGRAMS

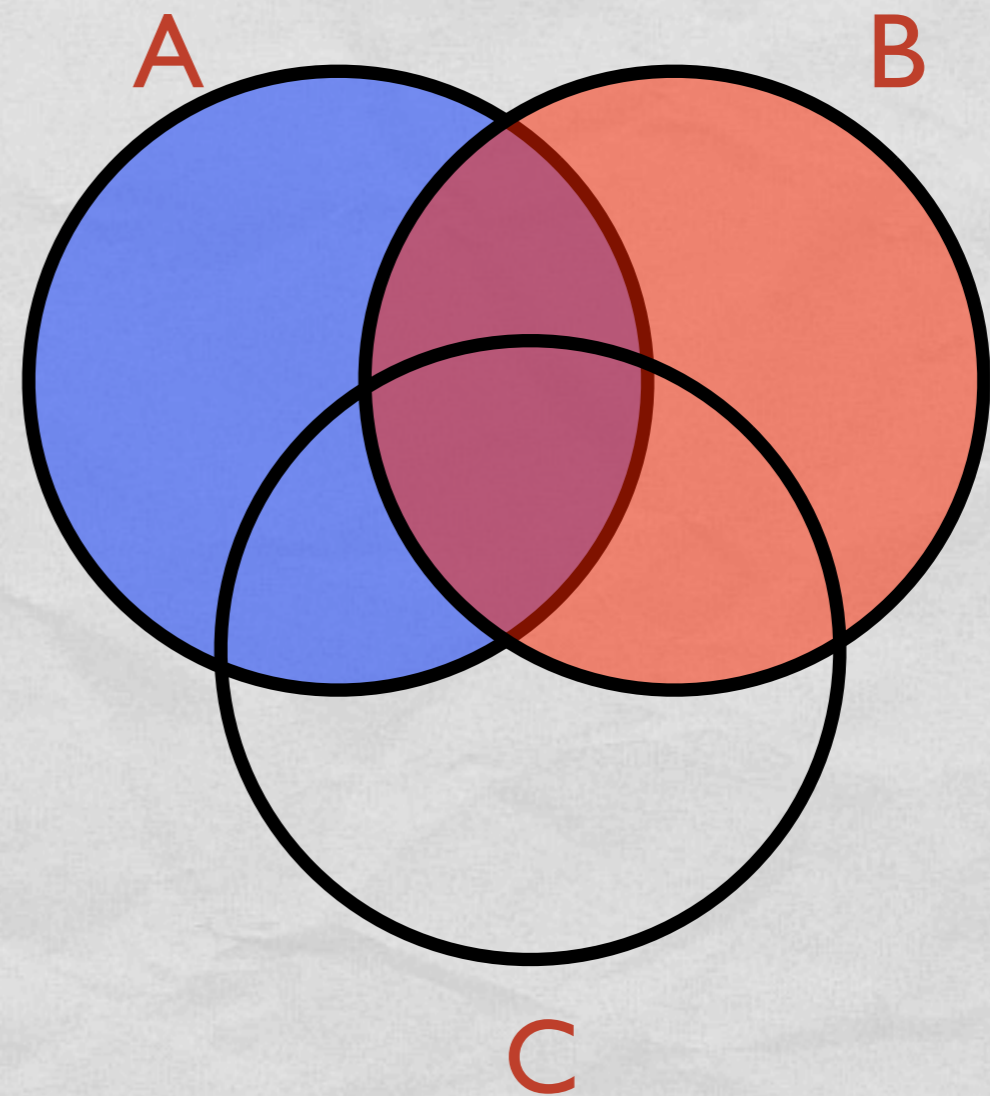
Everything in A = blue



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Everything in A = blue

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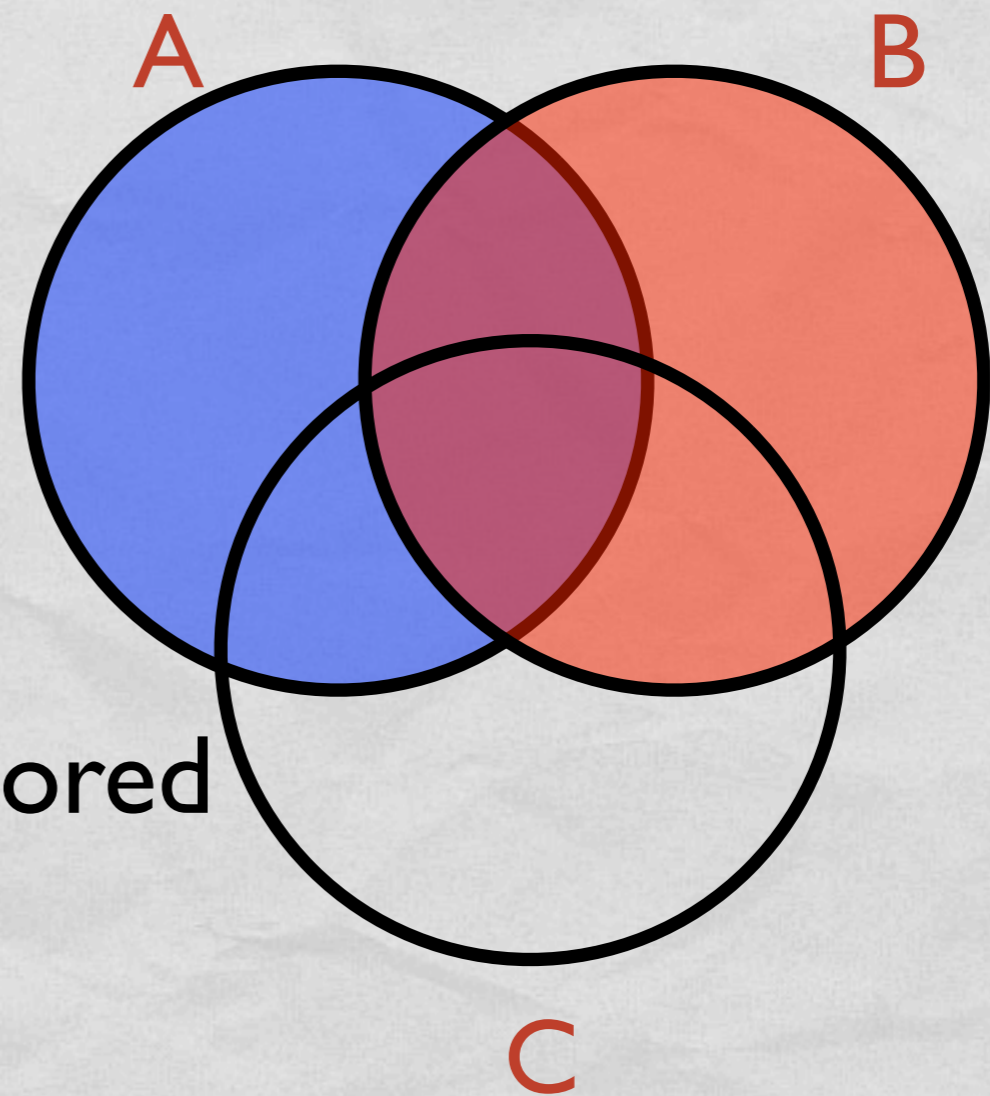
Everything in  $A$  = blue

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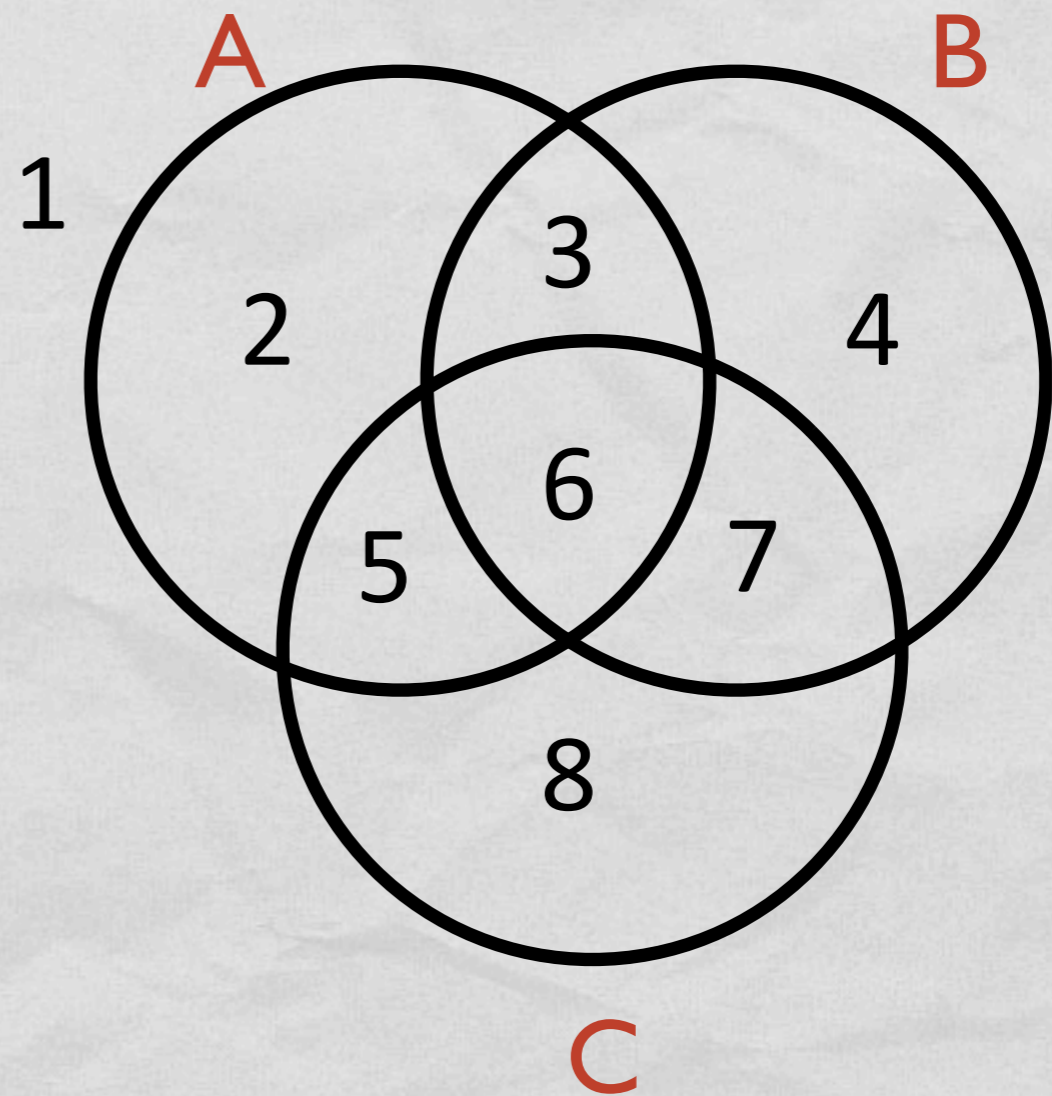
$A \cup B$  = colored

$A \cap B$  = purple

$C \setminus (A \cup B)$  = in  $C$  but not colored



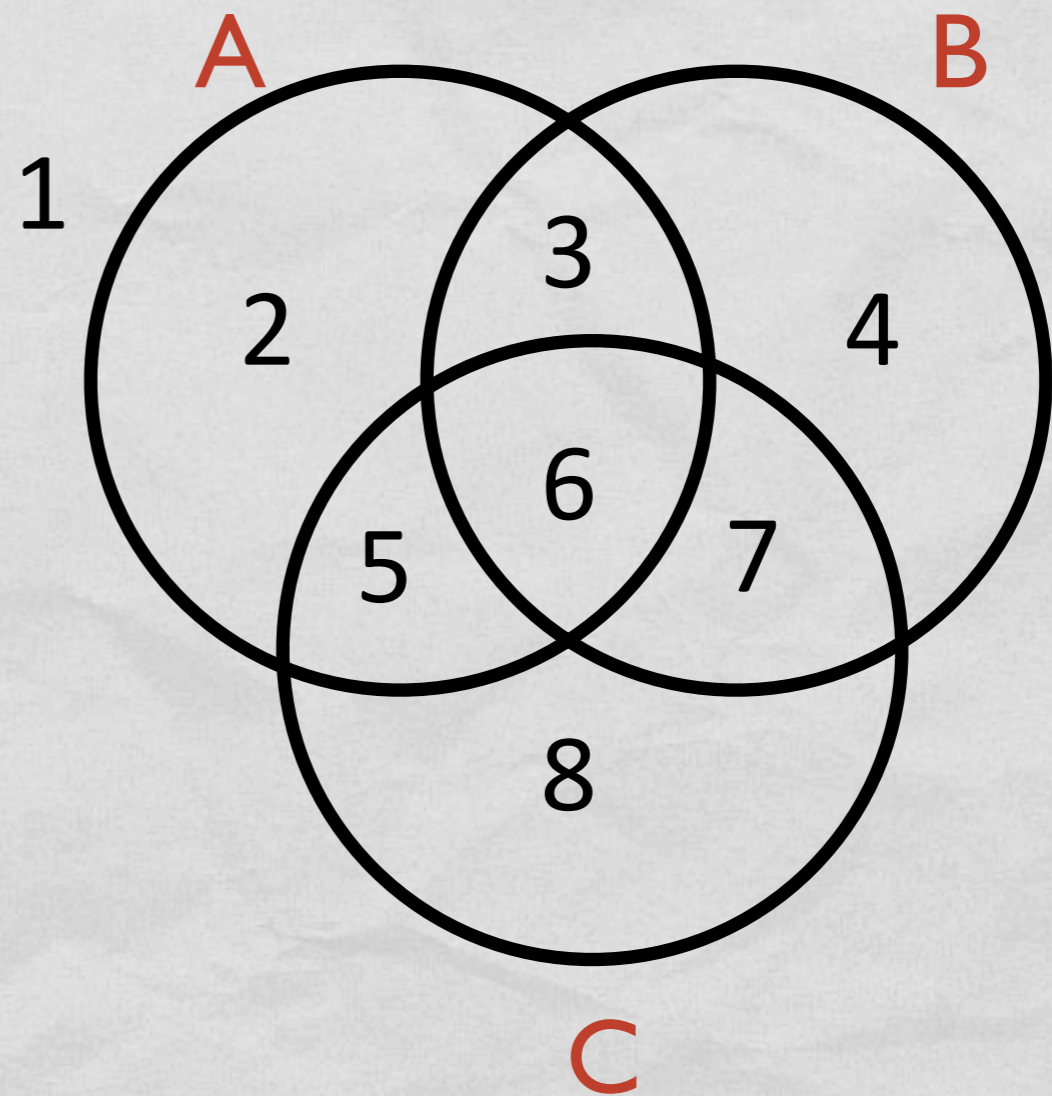
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region 6 =  $A \cap B \cap C$

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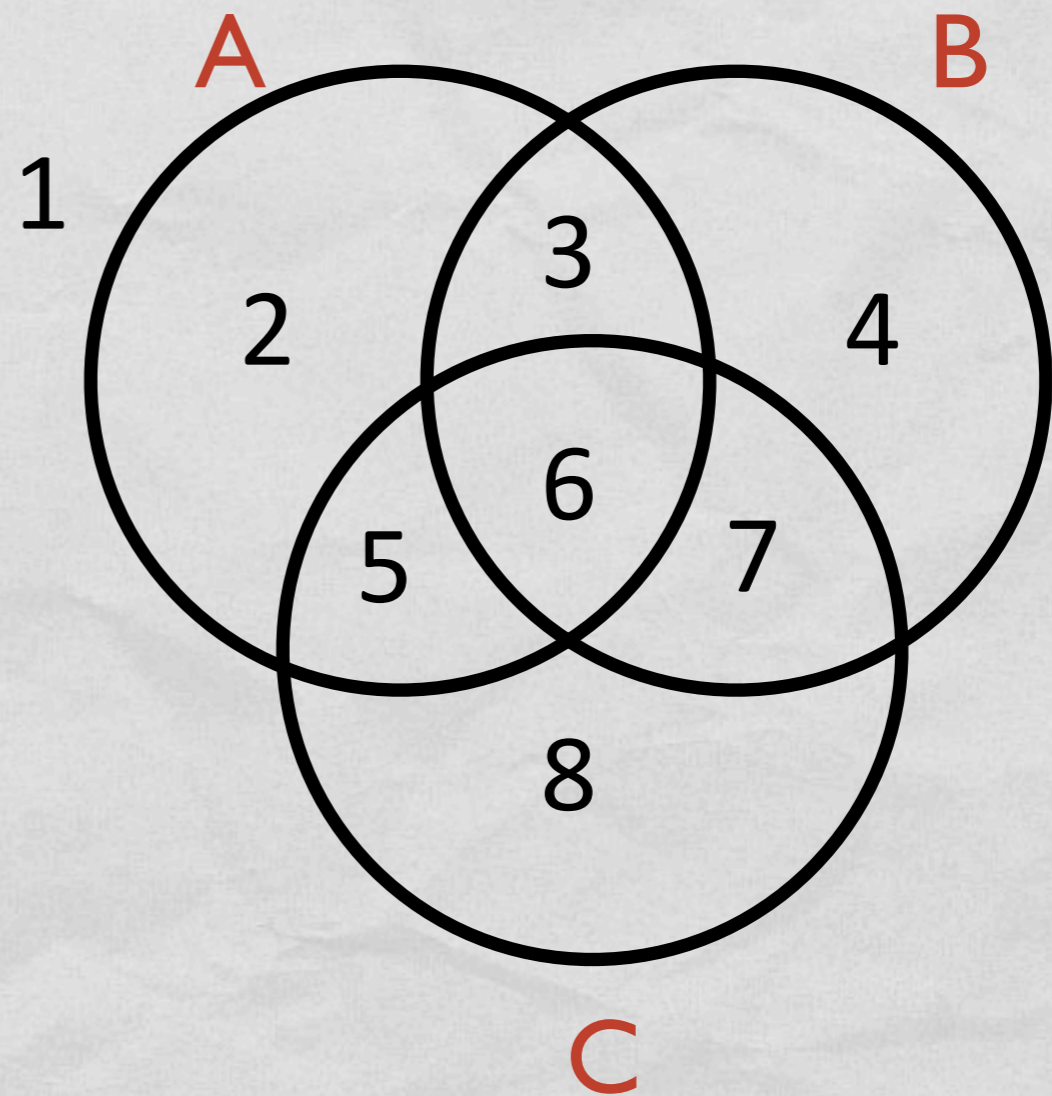
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region 6 =  $A \cap B \cap C$

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$$A = 2 + 3 + 5 + 6$$

$$\begin{aligned} A \cap B &= 2 + 3 + 5 + 6 \cap 3 + 4 + 6 + 7 \\ &= 3 + 6 \end{aligned}$$



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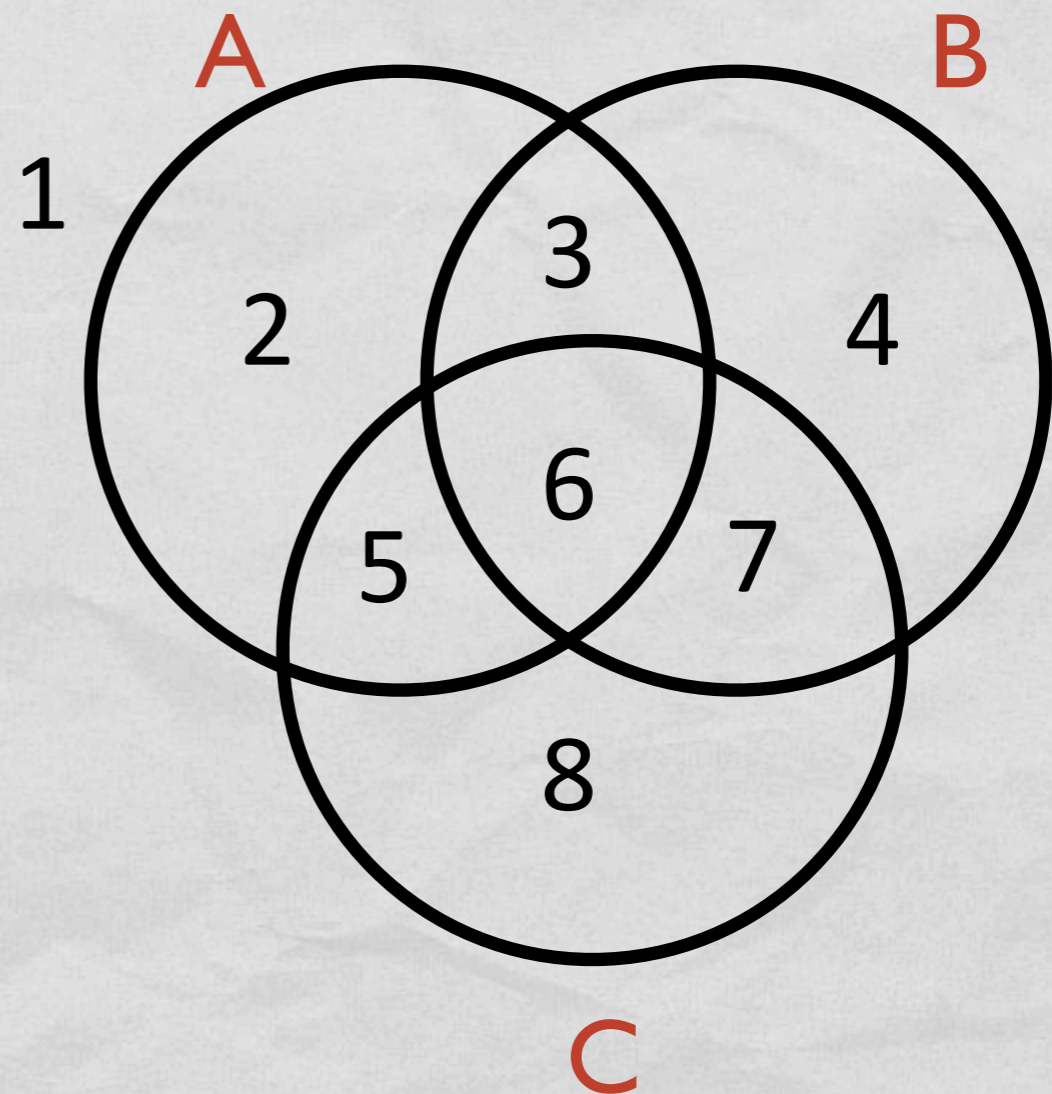
region 3 =  $A \cap B \cap C^c$

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region  $5 + 6 + 7 =$

$$C \cap (A \cup B) = (C \cap A) \cup (C \cap B)$$



# SUBSETS

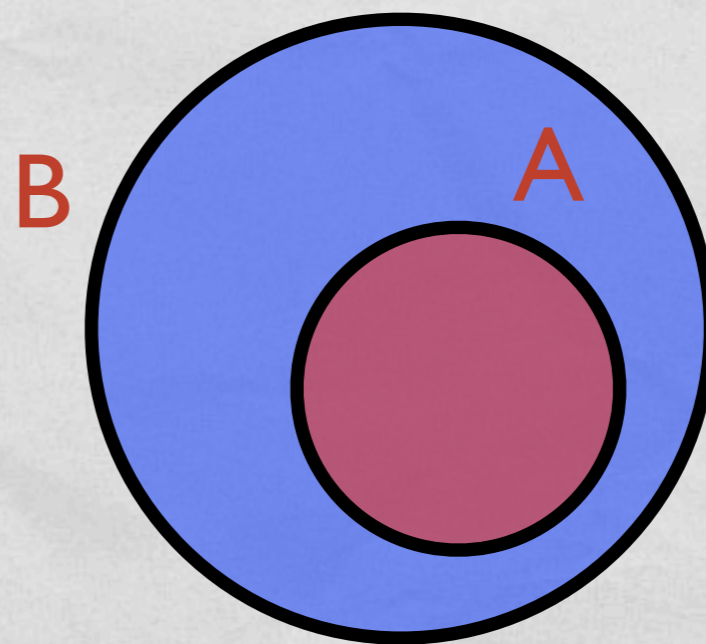


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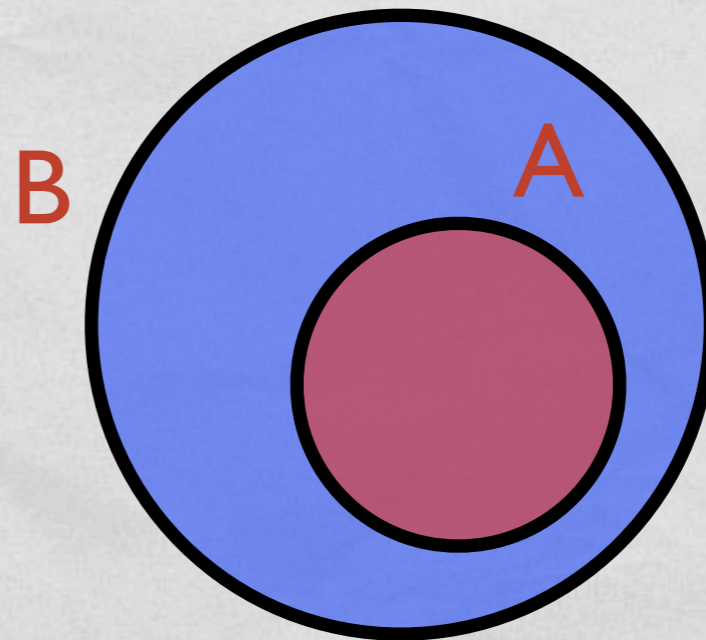
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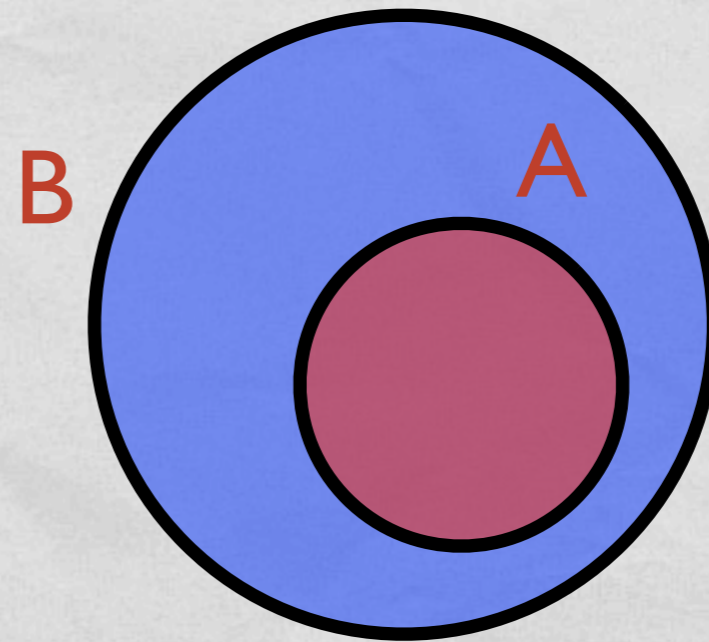
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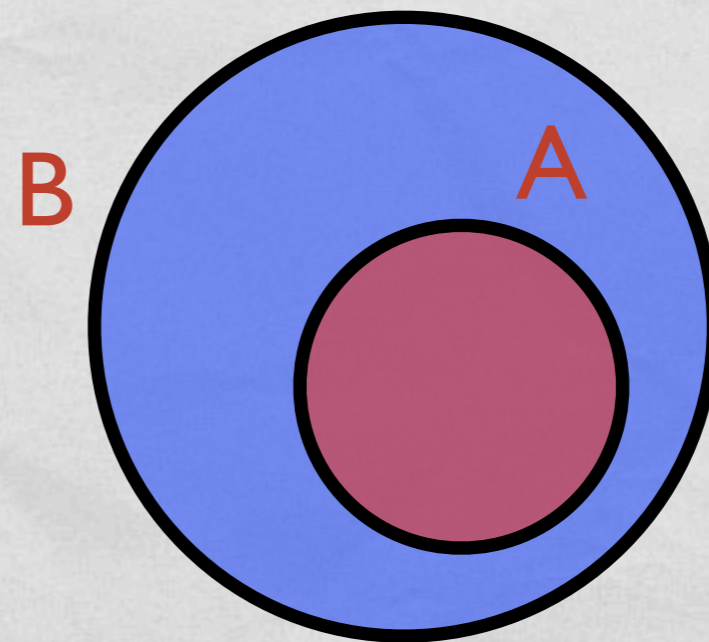
$$A \subsetneq B =_{\text{def}} A \subseteq B \wedge A \neq B$$

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$$\vdash (A \subseteq B \wedge B \subseteq A) \rightarrow A=B$$

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Membership isn't (generally)

$$(1 \in \{1,2\}, \{1,2\} \in \{\{1,2\}, \{2,3\}\})$$

$$\text{but } 1 \notin \{\{1,2\}, \{2,3\}\}$$



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and yes,  $\vdash \forall x (x < 2^x)$  even for infinite  $x$

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But we know that  $U \subseteq \wp(U)$  and  $\wp(U) \subseteq U$

From these two we can prove  $U = \wp(U)$

-- but we just showed that  $\vdash |\wp(U)| < |U|$

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$\vdash \neg \exists y \forall x [x \in y \leftrightarrow x \notin x]$  by Russell's argument

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So from a set  $z$ , you can shrink it down.

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+Replacement, Foundation, Choice make ZFC Set Theory