

In a certain place, all the inhabitants are either Knights or Knaves. Knights always tell the truth and Knaves never tell the truth.

You meet two inhabitants, A and B. A says "Both of us are knights." B says "A is a knave." What, if anything, can you infer from this?

# THE LOGIC OF ATOMIC SENTENCES: PROOFS OF (IN)VALIDITY

Friday, 24 January

# VALIDITY IN FOL

- A sentence S is a logical consequence of sentences
   P<sub>1</sub>...P<sub>n</sub> iff the argument with P<sub>1</sub>...P<sub>n</sub> as the premises and S as the conclusion is valid.
  - A formal deduction in  $\mathcal F$  proves validity.
- A sentence S is a <u>nonconsequence</u> of sentences P<sub>1</sub>...
   P<sub>n</sub> iff the argument with P<sub>1</sub>...P<sub>n</sub> as the premises and S as the conclusion is invalid.
  - A counterexample (such as a world in Tarski's World) proves invalidity.

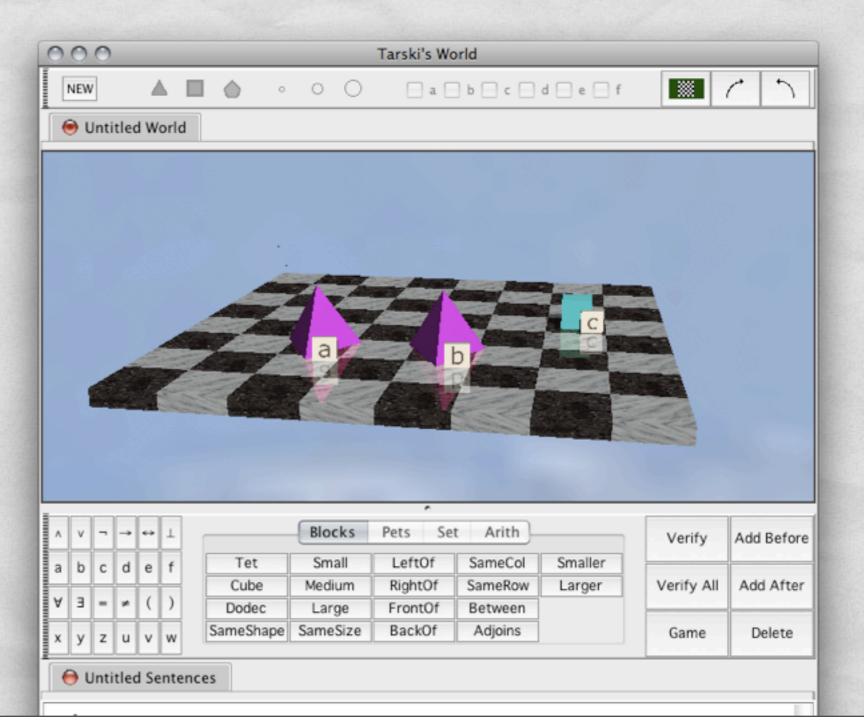
### Proving Nonconsequence

A Charles and a second state of the state

#### **Example:**

SameSize(a, b)
 Small(c)

3. Small(a)





#### **Example:**

I. LeftOf(a, b)2. LeftOf(a, c)3. b=d

4. LeftOf(d,c)

#### **Example:**

I. SameSize(a, b)2. SameSize(a, c)3. Medium(b)

4. Medium(c)

### Invalid

Valid

# HOW DOYOU KNOW?

- If you can find a counterexample, then you know the argument is invalid
- If you can't find a counterexample, you might need to keep looking
  - Or you could try to prove that the conclusion follows
  - If you do find a proof, you know it is valid

# FITCH-STYLE DEDUCTIVE SYSTEM

Contraction of the second second

I. P 2. Q 3. S<sub>1</sub> 4. S<sub>2</sub> .... n. S<sub>n</sub>

Justification 1 Justification 2

Justification n

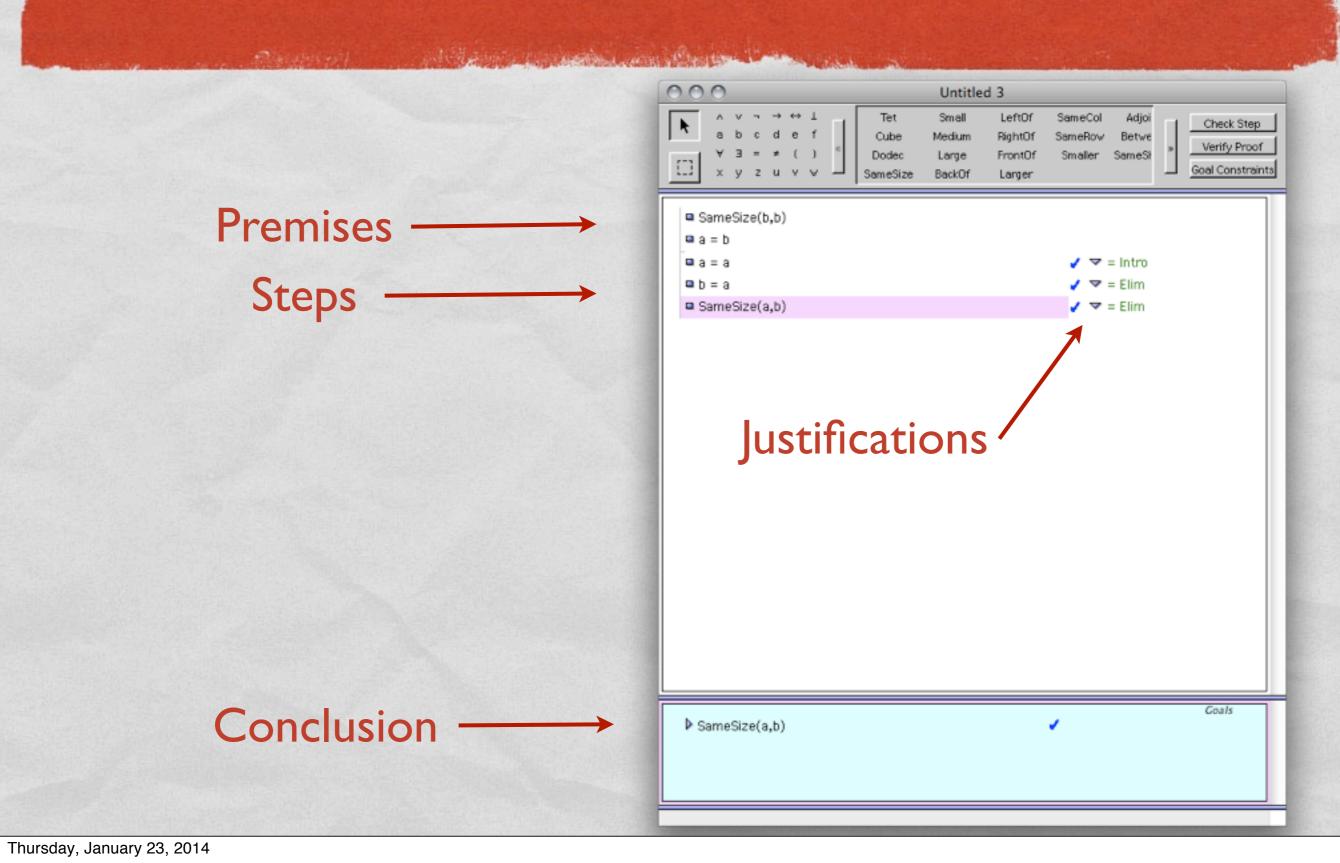
# FITCH-STYLE DEDUCTIVE SYSTEM

Rules of the system  $\mathcal{F}$ :

- Intro
- = Elim

Reit (Reiteration): "we have already shown that P"

## THE COMPUTER PROGRAM FITCH



### **PROOFS IN FITCH**

#### **Example:**

SameSize(a, b)
 a=c
 b=d

4. SameSize(c,d)

I. SameSize(a, b)Premise2. a=cPremise3. b=dPremise4. SameSize(c,b)= Elim 1,25. SameSize(c,d)= Elim 3,4

#### Valid

Thursday, January 23, 2014

### ANALYTICAL CONSEQUENCE

- There are no rules in  $\mathcal F$  that take advantage of the meaning of the blocks world predicates
  - LeftOf(a,b) is a logical consequence of RightOf(b,a) but you can't prove this in  $\mathcal{F}$
- The rule 'Analytical Consequence' (Ana Con) in Fitch makes use of these meanings

# ANALYTICAL CONSEQUENCE

#### Example:

SameSize(a, b)
 SameSize(a, c)
 Medium(b)

4. Medium(c)

Proof:

I. SameSize(a, b)
2. SameSize(a, c)
3. Medium(b)
4. SameSize(b,c) Ana Con 1,2
5. Medium(c) Ana Con 3,4

### Valid

Thursday, January 23, 2014