

PROOFS WITH CONDITIONALS

Monday, 24 February

CONDITIONAL ELIMINATION

- \rightarrow Elimination: from $P \rightarrow Q$ and P , we can infer Q .

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	1. $P \rightarrow Q$	
	2. P	
	—	
	3. Q	→ Elim: 1,2

CONDITIONAL ELIMINATION

- \rightarrow Elimination: from $P \rightarrow Q$ and P , we can infer Q .

	1. $P \rightarrow Q$	
	2. P	
	<hr/>	
	3. Q	\rightarrow Elim: 1,2

	1. $(A \vee B) \rightarrow (C \wedge D)$	
	2. $A \vee B$	
	<hr/>	
	3. $C \wedge D$	\rightarrow Elim: 1,2

EXAMPLE

Example:

	A	
	A	→ B
	B	→ C
	<hr/>	
	C	

EXAMPLE

Example:

A	
A → B	
B → C	
<hr/>	
C	

1. A	
2. A → B	
3. B → C	
<hr/>	

EXAMPLE

Example:

A
A \rightarrow B
B \rightarrow C
<hr/>
C

1. A
2. A \rightarrow B
3. B \rightarrow C
<hr/>
4. B

\rightarrow Elim 1,2

EXAMPLE

Example:

A	
A → B	
B → C	
<hr/>	
C	

1. A	
2. A → B	
3. B → C	
<hr/>	
4. B	→ Elim 1,2
5. C	→ Elim 3,4

EXAMPLE

Example:

$$\begin{array}{l} P \wedge Q \\ (P \vee R) \rightarrow S \\ (S \wedge Q) \rightarrow T \\ \hline T \end{array}$$

EXAMPLE

Example:

$$\begin{array}{l} P \wedge Q \\ (P \vee R) \rightarrow S \\ (S \wedge Q) \rightarrow T \\ \hline T \end{array}$$

$$\begin{array}{l} 1. P \wedge Q \\ 2. (P \vee R) \rightarrow S \\ 3. (S \wedge Q) \rightarrow T \\ \hline \end{array}$$

EXAMPLE

Example:

$P \wedge Q$
$(P \vee R) \rightarrow S$
$(S \wedge Q) \rightarrow T$
<hr/>
T

1. $P \wedge Q$	
2. $(P \vee R) \rightarrow S$	
3. $(S \wedge Q) \rightarrow T$	
<hr/>	
4. P	\wedge Elim 1

EXAMPLE

Example:

$P \wedge Q$
$(P \vee R) \rightarrow S$
$(S \wedge Q) \rightarrow T$
<hr/>
T

1. $P \wedge Q$	
2. $(P \vee R) \rightarrow S$	
3. $(S \wedge Q) \rightarrow T$	
<hr/>	
4. P	\wedge Elim 1
5. Q	\wedge Elim 1

EXAMPLE

Example:

$$\begin{array}{|l} P \wedge Q \\ (P \vee R) \rightarrow S \\ (S \wedge Q) \rightarrow T \\ \hline T \end{array}$$
$$\begin{array}{l} 1. P \wedge Q \\ 2. (P \vee R) \rightarrow S \\ 3. (S \wedge Q) \rightarrow T \\ \hline 4. P \quad \wedge \text{Elim 1} \\ 5. Q \quad \wedge \text{Elim 1} \\ 6. P \vee R \quad \vee \text{Intro 4} \end{array}$$

EXAMPLE

Example:

$P \wedge Q$
$(P \vee R) \rightarrow S$
$(S \wedge Q) \rightarrow T$
<hr/>
T

1. $P \wedge Q$	
2. $(P \vee R) \rightarrow S$	
3. $(S \wedge Q) \rightarrow T$	
<hr/>	
4. P	\wedge Elim 1
5. Q	\wedge Elim 1
6. $P \vee R$	\vee Intro 4
7. S	\rightarrow Elim 2,6

EXAMPLE

Example:

$P \wedge Q$
$(P \vee R) \rightarrow S$
$(S \wedge Q) \rightarrow T$
<hr/>
T

1. $P \wedge Q$	
2. $(P \vee R) \rightarrow S$	
3. $(S \wedge Q) \rightarrow T$	
<hr/>	
4. P	\wedge Elim 1
5. Q	\wedge Elim 1
6. $P \vee R$	\vee Intro 4
7. S	\rightarrow Elim 2,6
8. $S \wedge Q$	\wedge Intro 5,7

EXAMPLE

Example:

$P \wedge Q$
$(P \vee R) \rightarrow S$
$(S \wedge Q) \rightarrow T$
<hr/>
T

1. $P \wedge Q$	
2. $(P \vee R) \rightarrow S$	
3. $(S \wedge Q) \rightarrow T$	
<hr/>	
4. P	\wedge Elim 1
5. Q	\wedge Elim 1
6. $P \vee R$	\vee Intro 4
7. S	\rightarrow Elim 2,6
8. $S \wedge Q$	\wedge Intro 5,7
9. T	\rightarrow Elim 3,8

RULES FOR CONDITIONALS

RULES FOR CONDITIONALS

1. $P \rightarrow Q$

2. P

3. Q \rightarrow Elim: 1,2

RULES FOR CONDITIONALS

1. $P \rightarrow Q$
2. P
—
3. Q \rightarrow Elim: 1,2

1. $P \rightarrow Q$
2. $\neg P$
—
3. $\neg Q$ **INVALID**

RULES FOR CONDITIONALS

1. $P \rightarrow Q$
2. P
—
3. Q \rightarrow Elim: 1,2

1. $P \rightarrow Q$
2. $\neg P$
—
3. $\neg Q$ **INVALID**

1. $P \rightarrow Q$
2. Q
—
3. P **INVALID**

RULES FOR CONDITIONALS

1. $P \rightarrow Q$
2. P
—
3. Q \rightarrow Elim: 1,2

1. $P \rightarrow Q$
2. $\neg P$
—
3. $\neg Q$ **INVALID**

1. $P \rightarrow Q$
2. Q
—
3. P **INVALID**

1. $P \rightarrow Q$
2. $\neg Q$
—
3. $\neg P$ **VALID, but not \rightarrow E**

RULES FOR CONDITIONALS

1. $P \rightarrow Q$	
2. P	
<hr/>	
3. Q	\rightarrow Elim: 1,2

Modus Ponens

1. $P \rightarrow Q$	
2. $\neg Q$	
<hr/>	
3. $\neg P$	VALID, but not $\rightarrow E$

Modus Tollens

MODUS TOLLENS

Example:

1. $P \rightarrow Q$

2. $\neg Q$

3. $\neg P$

MODUS TOLLENS

Example:

1. $P \rightarrow Q$
2. $\neg Q$
—
3. $\neg P$

1. $P \rightarrow Q$
2. $\neg Q$
—
 $\neg P$

MODUS TOLLENS

Example:

1. $P \rightarrow Q$
2. $\neg Q$
—
3. $\neg P$

1. $P \rightarrow Q$
2. $\neg Q$
—
3. P for \neg Intro
—
 $\neg P$ \neg Intro

MODUS TOLLENS

Example:

1. $P \rightarrow Q$
2. $\neg Q$
—
3. $\neg P$

1. $P \rightarrow Q$
2. $\neg Q$
—
3. P for \neg Intro
4. Q \rightarrow Elim 1,3
—
 $\neg P$ \neg Intro

MODUS TOLLENS

Example:

	1. $P \rightarrow Q$
	2. $\neg Q$

	3. $\neg P$

	1. $P \rightarrow Q$		
	2. $\neg Q$		

		3. P	for \neg Intro

		4. Q	\rightarrow Elim 1,3
		5. \perp	\perp Intro 2,4

		$\neg P$	\neg Intro

MODUS TOLLENS

Example:

	1. $P \rightarrow Q$
	2. $\neg Q$

	3. $\neg P$

	1. $P \rightarrow Q$		
	2. $\neg Q$		

		3. P	for \neg Intro
		4. Q	\rightarrow Elim 1,3
		5. \perp	\perp Intro 2,4
		6. $\neg P$	\neg Intro 3-5

BICONDITIONAL ELIMINATION

- \leftrightarrow Elimination: from $P \leftrightarrow Q$ and P/Q , we can infer Q/P .

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- \leftrightarrow Elimination: from $P \leftrightarrow Q$ and P/Q , we can infer Q/P .

1. $P \leftrightarrow Q$	
2. P	

3. Q	\leftrightarrow Elim: 1,2

BICONDITIONAL ELIMINATION

- \leftrightarrow Elimination: from $P \leftrightarrow Q$ and P/Q , we can infer Q/P .

	1. $P \leftrightarrow Q$	
	2. P	
	—	
	3. Q	\leftrightarrow Elim: 1,2

	1. $P \leftrightarrow Q$	
	2. Q	
	—	
	3. P	\leftrightarrow Elim: 1,2

EXAMPLE

Example:

$$\begin{array}{l} A \\ A \leftrightarrow B \\ \neg C \leftrightarrow B \\ \hline \neg C \end{array}$$

EXAMPLE

Example:

	A
	$A \leftrightarrow B$
	$\neg C \leftrightarrow B$
	—
	$\neg C$

	1. A
	2. $A \leftrightarrow B$
	3. $\neg C \leftrightarrow B$
	—

EXAMPLE

Example:

	A
	$A \leftrightarrow B$
	$\neg C \leftrightarrow B$
	—
	$\neg C$

	1. A
	2. $A \leftrightarrow B$
	3. $\neg C \leftrightarrow B$
	—
	4. B
	\leftrightarrow Elim 1,2

EXAMPLE

Example:

	A	
	$A \leftrightarrow B$	
	$\neg C \leftrightarrow B$	
	—	
	$\neg C$	

	1. A	
	2. $A \leftrightarrow B$	
	3. $\neg C \leftrightarrow B$	
	—	
	4. B	\leftrightarrow Elim 1,2
	5. $\neg C$	\leftrightarrow Elim 3,4

BACK TO CASE ARGUMENTS

Example:

$$\begin{array}{|l} P \vee Q \\ P \rightarrow R \\ Q \rightarrow R \\ \hline R \end{array}$$

BACK TO CASE ARGUMENTS

Example:

$$\begin{array}{|l} P \vee Q \\ P \rightarrow R \\ Q \rightarrow R \\ \hline R \end{array}$$
$$\begin{array}{|l} 1. P \vee Q \\ 2. P \rightarrow R \\ 3. Q \rightarrow R \\ \hline \end{array}$$
 R $\vee\text{Elim } 1, 4\text{-}\dots$

BACK TO CASE ARGUMENTS

Example:

$$\begin{array}{|l} P \vee Q \\ P \rightarrow R \\ Q \rightarrow R \\ \hline R \end{array}$$
$$\begin{array}{|l} 1. P \vee Q \\ 2. P \rightarrow R \\ 3. Q \rightarrow R \\ \hline \begin{array}{|l} 4. P \quad \text{for } \vee\text{Elim} \\ \hline R \\ \end{array} \\ \begin{array}{|l} Q \quad \text{for } \vee\text{Elim} \\ \hline R \\ \end{array} \\ R \quad \vee\text{Elim } 1, 4-... \end{array}$$

BACK TO CASE ARGUMENTS

Example:

	$P \vee Q$
	$P \rightarrow R$
	$Q \rightarrow R$

	R

	1. $P \vee Q$	
	2. $P \rightarrow R$	
	3. $Q \rightarrow R$	

	4. P	for \vee Elim

	5. R	\rightarrow Elim 2,4
	6. Q	for \vee Elim

	7. R	\rightarrow Elim 3,6
	8. R	\vee Elim 1, 4-5, 6-7

BACK TO CASE ARGUMENTS

Example:

$P \vee Q$
$P \rightarrow R$
$Q \rightarrow R$
—
R

from P ,
infer R

Proof from
 P to R

1. $P \vee Q$	
2. $P \rightarrow R$	
3. $Q \rightarrow R$	
—	
4. P	for \vee Elim
—	
5. R	\rightarrow Elim 2,4
6. Q	for \vee Elim
—	
7. R	\rightarrow Elim 3,6
8. R	\vee Elim 1, 4-5, 6-7

SUBPROOFS AND CONDITIONALS

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- The conditional $P \rightarrow R$ says that if P is true, then R is also true

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- A subproof from P to R shows that if P were true, then R would also be true.

SUBPROOFS AND CONDITIONALS

- The conditional $P \rightarrow R$ says that if P is true, then R is also true
- A subproof from P to R shows that if P were true, then R would also be true.
- In \mathcal{F} , the disjunction elimination rule (\vee Elim - case argument) uses a disjunction and two subproofs. In many other systems (and on Wikipedia for example) you need the disjunction and two conditionals

\vee Elim (alternate version)

Example:

1.	$P \vee Q$	
2.	$P \rightarrow R$	
3.	$Q \rightarrow R$	
<hr/>		
4.	R	\vee Elim 1, 2, 3

\vee Elim (alternate version)

Example:

1.	$P \vee Q$	
2.	$P \rightarrow R$	
3.	$Q \rightarrow R$	
<hr/>		
4.	R	\vee Elim 1, 2, 3

The missing subproofs just are the conditionals

SUBPROOFS AND CONDITIONALS

- The conditional $P \rightarrow R$ says that if P is true, then R is also true
- A subproof from P to R shows that if P were true, then R would also be true.
- The Introduction rule for the conditional says that to prove that a conditional $P \rightarrow R$ is true, assume P and show that it does lead to R

CONDITIONAL INTRODUCTION

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- \rightarrow Introduction

From a subproof from P to Q , we can infer $P \rightarrow Q$.

CONDITIONAL INTRODUCTION

- \rightarrow Introduction

From a subproof from P to Q , we can infer $P \rightarrow Q$.



CONDITIONAL INTRODUCTION

- \rightarrow Introduction

From a subproof from P to Q , we can infer $P \rightarrow Q$.



This rule is often known as Conditional Proof

CHAIN ARGUMENT

Example:

$$\begin{array}{l} P \rightarrow Q \\ Q \rightarrow R \\ \hline P \rightarrow R \end{array}$$

CHAIN ARGUMENT

Example:

$$\begin{array}{l} P \rightarrow Q \\ Q \rightarrow R \\ \hline P \rightarrow R \end{array}$$

$$\begin{array}{l} 1. P \rightarrow Q \\ 2. Q \rightarrow R \\ \hline P \rightarrow R \end{array}$$

CHAIN ARGUMENT

Example:

$$\begin{array}{l} P \rightarrow Q \\ Q \rightarrow R \\ \hline P \rightarrow R \end{array}$$

$$\begin{array}{l} 1. P \rightarrow Q \\ 2. Q \rightarrow R \\ \hline 3. P \\ \hline R \\ \hline P \rightarrow R \end{array}$$

for \rightarrow Intro

\rightarrow Intro

CHAIN ARGUMENT

Example:

$$\begin{array}{l} P \rightarrow Q \\ Q \rightarrow R \\ \hline P \rightarrow R \end{array}$$
$$\begin{array}{l} 1. P \rightarrow Q \\ 2. Q \rightarrow R \\ \hline \begin{array}{l} 3. P \\ \hline 4. Q \\ R \end{array} \\ P \rightarrow R \end{array}$$

for \rightarrow Intro

\rightarrow Elim 1,3

\rightarrow Intro

CHAIN ARGUMENT

Example:

$$\begin{array}{l} P \rightarrow Q \\ Q \rightarrow R \\ \hline P \rightarrow R \end{array}$$
$$\begin{array}{l} 1. P \rightarrow Q \\ 2. Q \rightarrow R \\ \hline \end{array}$$
$$\begin{array}{l} 3. P \\ \hline \end{array}$$
$$4. Q$$
$$5. R$$
$$P \rightarrow R$$

for \rightarrow Intro

\rightarrow Elim 1,3

\rightarrow Elim 2,4

\rightarrow Intro

CHAIN ARGUMENT

Example:

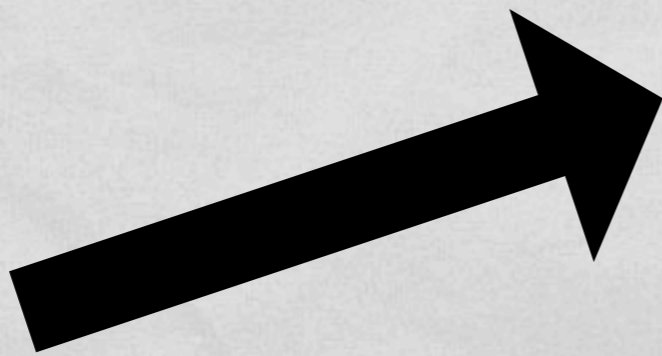
$$\begin{array}{l} P \rightarrow Q \\ Q \rightarrow R \\ \hline P \rightarrow R \end{array}$$
$$\begin{array}{l} 1. P \rightarrow Q \\ 2. Q \rightarrow R \\ \hline 3. P \quad \text{for } \rightarrow \text{Intro} \\ 4. Q \quad \rightarrow \text{Elim } 1,3 \\ 5. R \quad \rightarrow \text{Elim } 2,4 \\ 6. P \rightarrow R \quad \rightarrow \text{Intro } 3-5 \end{array}$$

CHAIN ARGUMENT

Example:

$$\begin{array}{|l} P \rightarrow Q \\ Q \rightarrow R \\ \hline P \rightarrow R \end{array}$$

A proof
from P to R


$$\begin{array}{|l} 1. P \rightarrow Q \\ 2. Q \rightarrow R \\ \hline 3. P \\ \hline 4. Q \\ 5. R \\ \hline 6. P \rightarrow R \end{array}$$

for \rightarrow Intro
 \rightarrow Elim 1,3
 \rightarrow Elim 2,4
 \rightarrow Intro 3-5

CHAIN ARGUMENT

Example:

$$\begin{array}{|l} P \rightarrow Q \\ Q \rightarrow R \\ \hline P \rightarrow R \end{array}$$

A proof
from P to R

So P leads to R

$$\begin{array}{|l} 1. P \rightarrow Q \\ 2. Q \rightarrow R \\ \hline 3. P \\ \hline 4. Q \\ 5. R \\ \hline 6. P \rightarrow R \end{array}$$

for \rightarrow Intro
 \rightarrow Elim 1,3
 \rightarrow Elim 2,4
 \rightarrow Intro 3-5

IMPORT/EXPORT LAWS

Example:

$$\begin{array}{|l} P \rightarrow (Q \rightarrow R) \\ \hline (P \wedge Q) \rightarrow R \end{array}$$

IMPORT/EXPORT LAWS

Example:

$$\frac{P \rightarrow (Q \rightarrow R)}{(P \wedge Q) \rightarrow R}$$

$$\frac{\frac{1. P \rightarrow (Q \rightarrow R)}{\frac{2. P \wedge Q}{R} \text{ for } \rightarrow\text{Intro}}{(P \wedge Q) \rightarrow R} \rightarrow\text{Intro}}$$

IMPORT/EXPORT LAWS

Example:

$$\frac{P \rightarrow (Q \rightarrow R)}{(P \wedge Q) \rightarrow R}$$

$$\begin{array}{l} 1. P \rightarrow (Q \rightarrow R) \\ \hline 2. P \wedge Q \quad \text{for } \rightarrow\text{Intro} \\ \hline 3. P \quad \wedge \text{Elim } 2 \\ \hline R \\ \hline (P \wedge Q) \rightarrow R \quad \rightarrow\text{Intro} \end{array}$$

IMPORT/EXPORT LAWS

Example:

$$\frac{P \rightarrow (Q \rightarrow R)}{(P \wedge Q) \rightarrow R}$$

$$\begin{array}{l} 1. P \rightarrow (Q \rightarrow R) \\ \hline 2. P \wedge Q \quad \text{for } \rightarrow\text{Intro} \\ \hline 3. P \quad \wedge \text{Elim } 2 \\ 4. Q \quad \wedge \text{Elim } 2 \\ \hline R \\ \hline (P \wedge Q) \rightarrow R \quad \rightarrow\text{Intro} \end{array}$$

IMPORT/EXPORT LAWS

Example:

$$\frac{P \rightarrow (Q \rightarrow R)}{(P \wedge Q) \rightarrow R}$$

$$\begin{array}{l} 1. P \rightarrow (Q \rightarrow R) \\ \hline 2. P \wedge Q \quad \text{for } \rightarrow\text{Intro} \\ \hline 3. P \quad \wedge \text{Elim } 2 \\ 4. Q \quad \wedge \text{Elim } 2 \\ 5. Q \rightarrow R \quad \rightarrow \text{Elim } 1,3 \\ \hline R \\ \hline (P \wedge Q) \rightarrow R \quad \rightarrow\text{Intro} \end{array}$$

IMPORT/EXPORT LAWS

Example:

$$\frac{P \rightarrow (Q \rightarrow R)}{(P \wedge Q) \rightarrow R}$$

$$\begin{array}{l} 1. P \rightarrow (Q \rightarrow R) \\ \hline 2. P \wedge Q \quad \text{for } \rightarrow \text{Intro} \\ \hline 3. P \quad \wedge \text{Elim } 2 \\ 4. Q \quad \wedge \text{Elim } 2 \\ 5. Q \rightarrow R \quad \rightarrow \text{Elim } 1,3 \\ 6. R \quad \rightarrow \text{Elim } 4,5 \\ \hline 7. (P \wedge Q) \rightarrow R \quad \rightarrow \text{Intro } 2-6 \end{array}$$